



Efficient Critical Paths Search Algorithm using Mergeable Heap

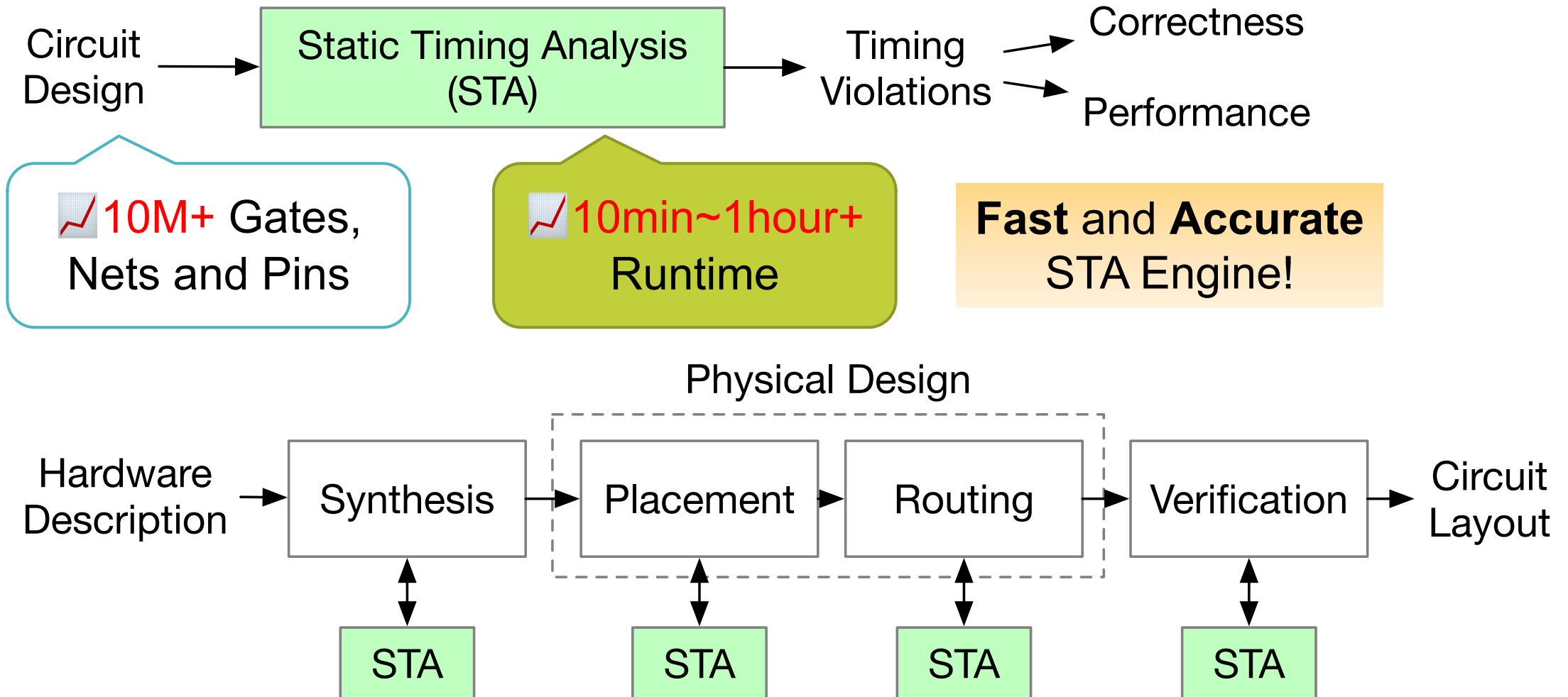
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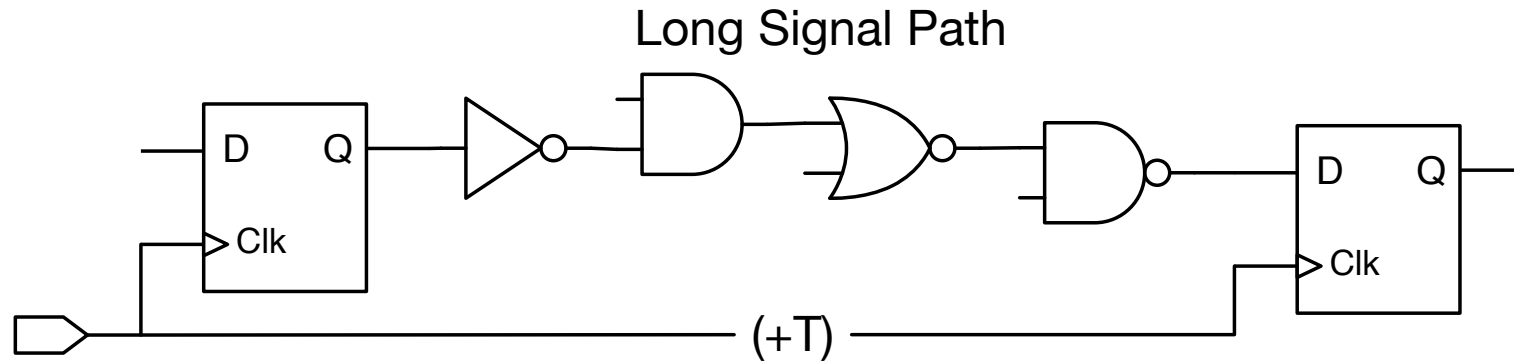
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Static Timing Analysis (STA)



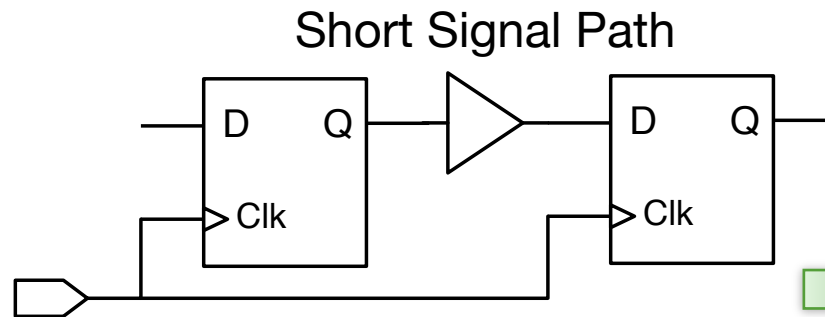
Critical Paths Searching in STA

Setup
Violation



➔ Top-*k* longest paths

Hold
Violation



➔ Top-*k* shortest paths

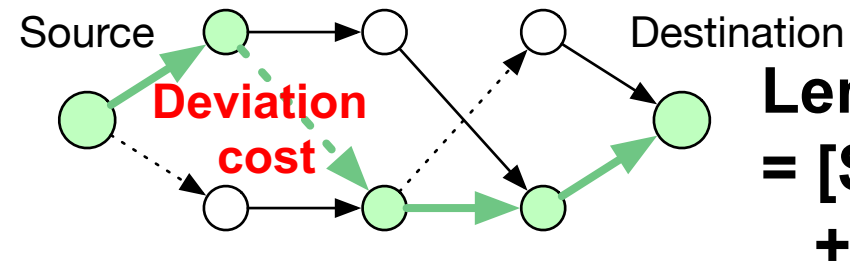
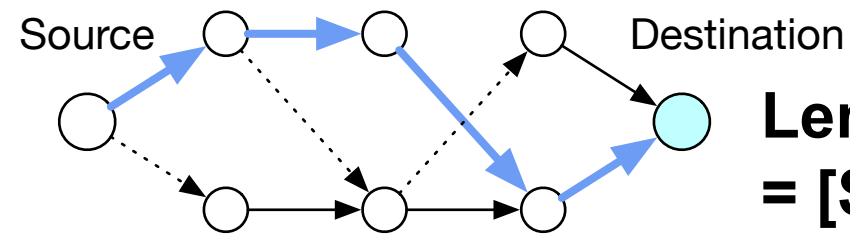
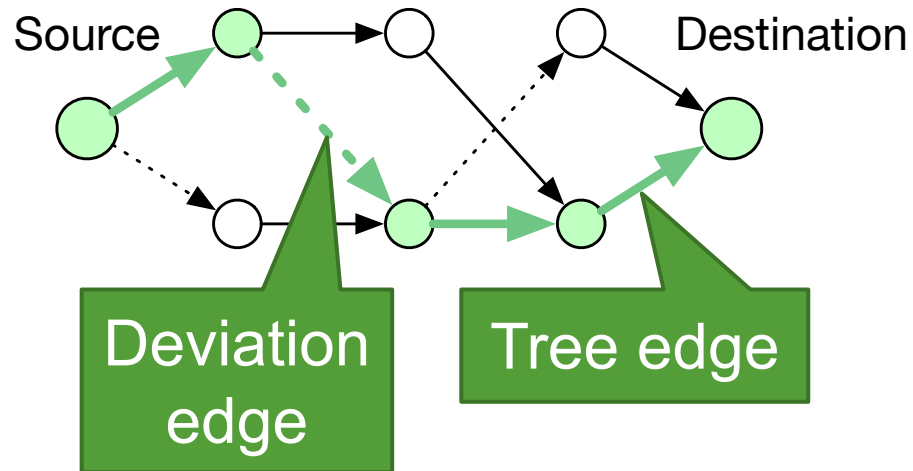
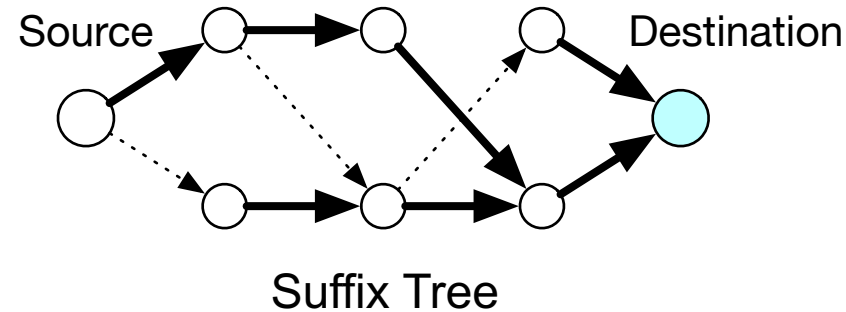
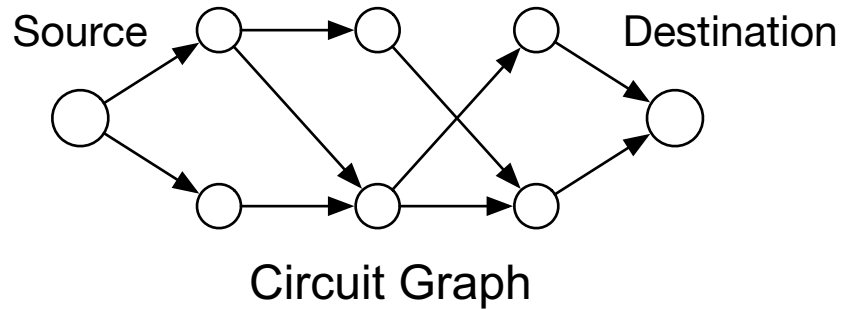
Input: Circuit graph, *k*

Output: Top-*k* shortest paths

Real case STA: *k* = 10,000 ~ 100,000

The State-of-the-art K -Shortest Path Algorithm

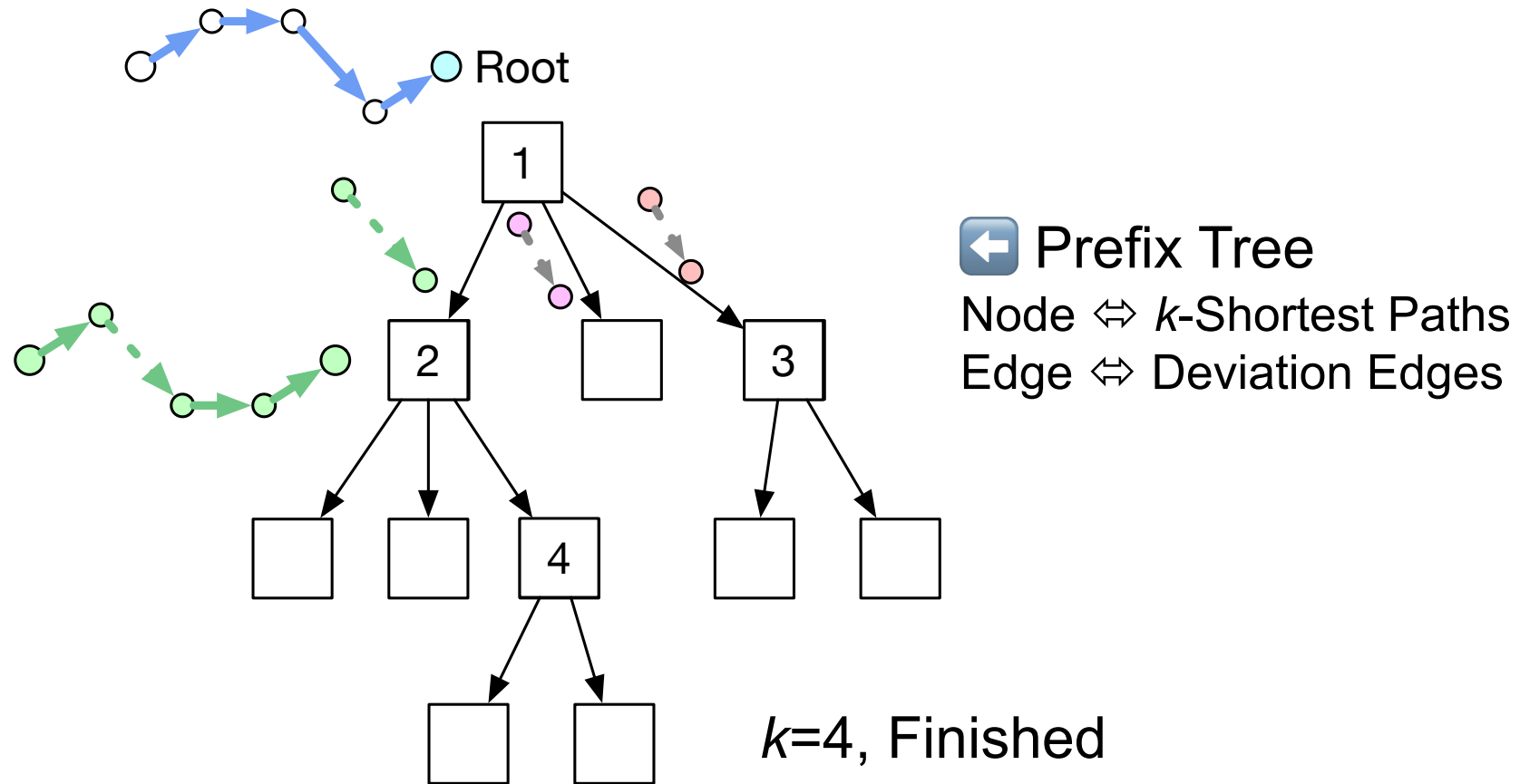
Suffix Tree - Prefix Tree Algorithm [OpenTimer, TCAD'20] [Guo, ICCAD'21]



Path \Leftrightarrow The set of deviation edges on the path { $\text{green circle} \dashrightarrow \text{green circle}$, ... }

The State-of-the-art K -Shortest Path Algorithm

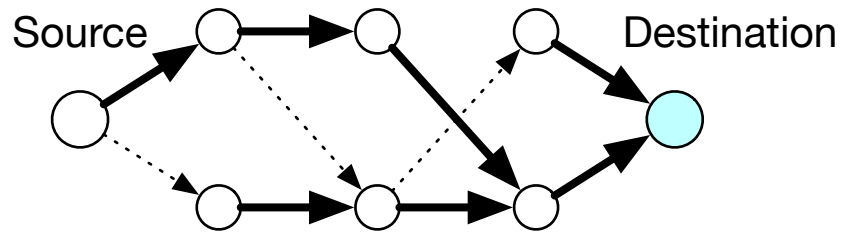
Suffix Tree - Prefix Tree Algorithm [OpenTimer, TCAD'20] [Guo, ICCAD'21]



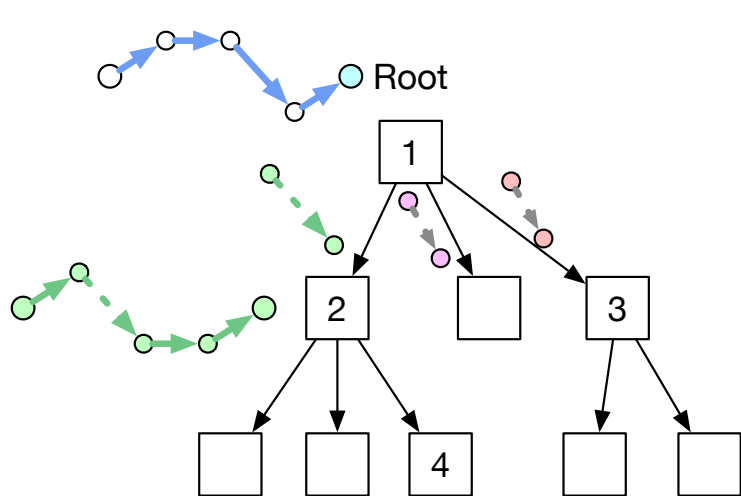
The Burden for Suffix-Prefix Algorithm

➤ $n = \#Pins + \#Edges$, $k = \#Paths$

➤ Step 1: Build the suffix tree

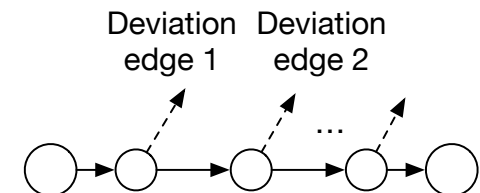


➤ Step 2: Search for top- k paths



$$k * \begin{array}{c} \boxed{2} \\ \swarrow \quad \downarrow \quad \searrow \\ \square \quad \square \quad \boxed{4} \end{array} = O(kn \log n)$$

Each exploration:
at most $O(n)$ deviation edges



Time
Complexity

$O(n)$
on acyclic graph

Expected:
 $\sim O(k+n)$

↑ This work

Total:

$O(kn \log n)$

Large?

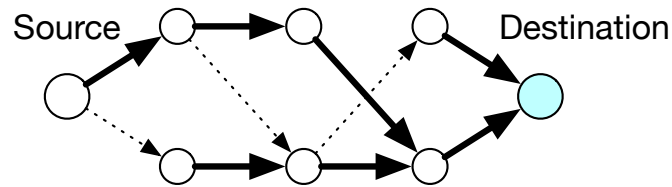
Large?

Our Contributions

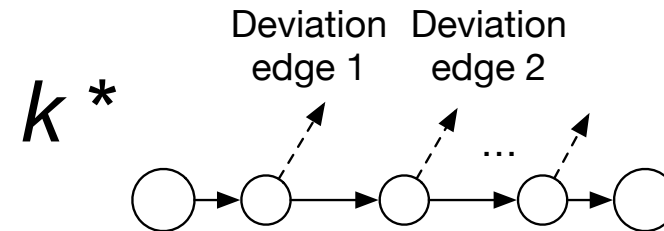
- A novel k -shortest path searching algorithm that runs in $O(n \log n + k \log k)$, asymptotically lower than baselines $O(kn \log n)$.
- Incorporating **persistent mergeable heaps** to store all path deviations for fast merging and duplicating.
- Introducing a novel **deviation preprocessing** step to precompute path deviations and speed up path searching.

Motivation: Pre-computing for Future Use

Step 1: Suffix tree



Step 2: Prefix tree



Each exploration: at most $O(n)$ deviation edges

Baseline

$O(n)$

$= O(kn \log n)$

Our
Algorithm

Prepare
deviation
edges

$O(n \log n)$

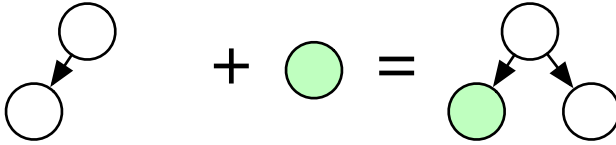
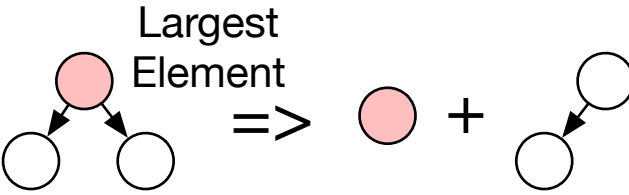
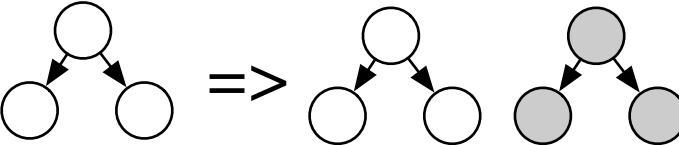
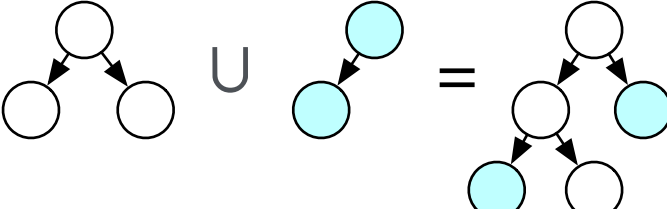
Use
precomputed
edges

$O(k \log k)$

Overall:

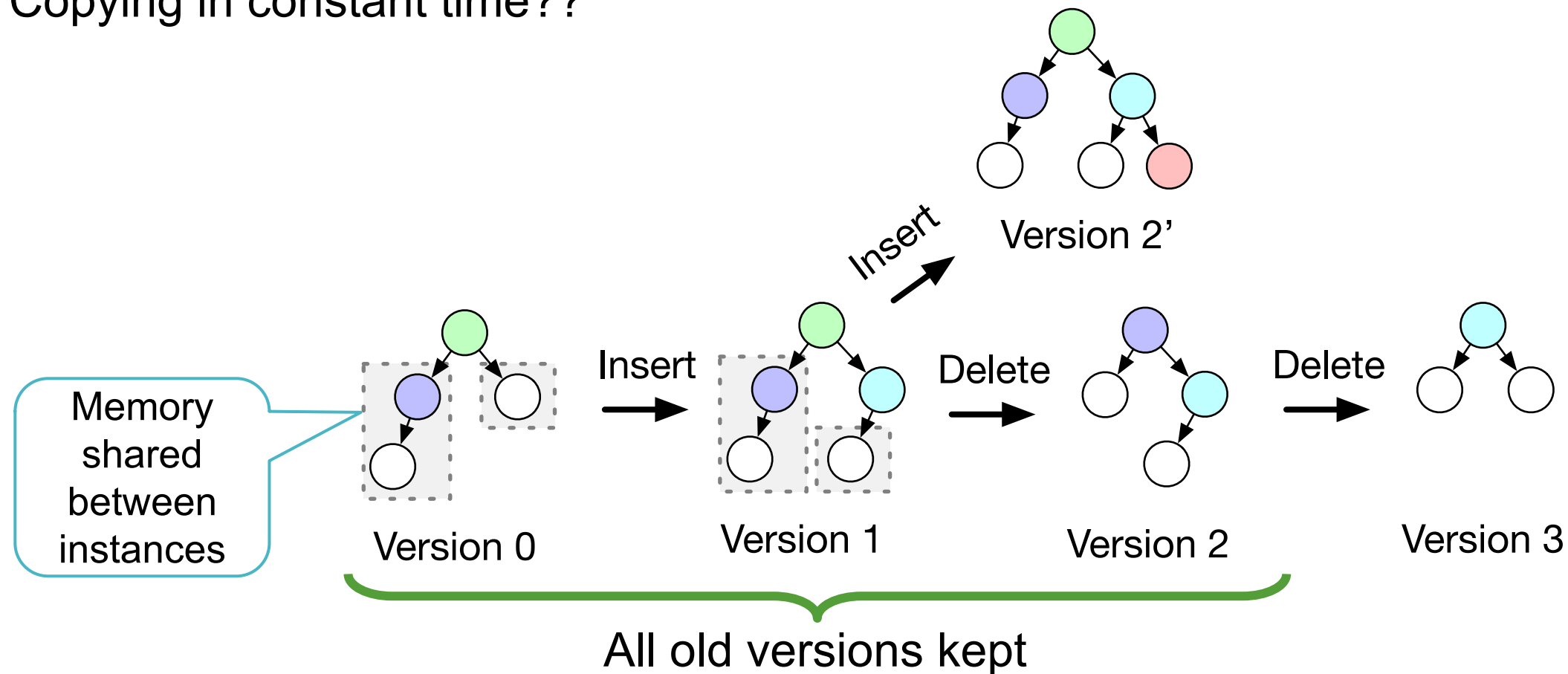
$O(n \log n + k \log k)$

The Building Block: Persistent Mergeable Heaps

	Ordinary Heaps (e.g. Binary heap)	Persistent Mergeable Heaps (e.g. Leftist tree)
	Insertion	$O(\log n)$
 <p>Largest Element</p>	Deletion (largest element)	$O(\log n)$
	Copying	$O(n)$
	Merging	$O(n)$

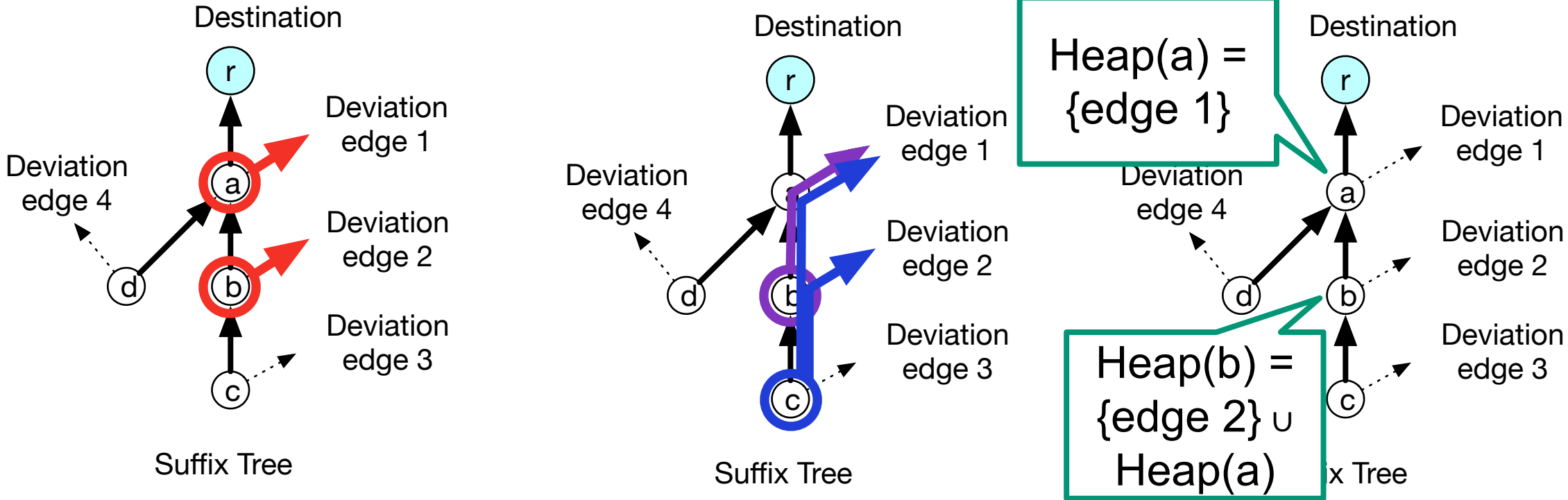
Magic Behind Leftist-Tree: **Persistency**

Copying in constant time??



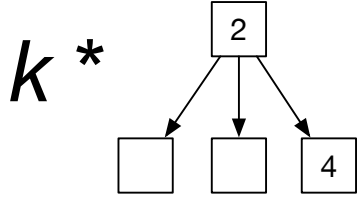
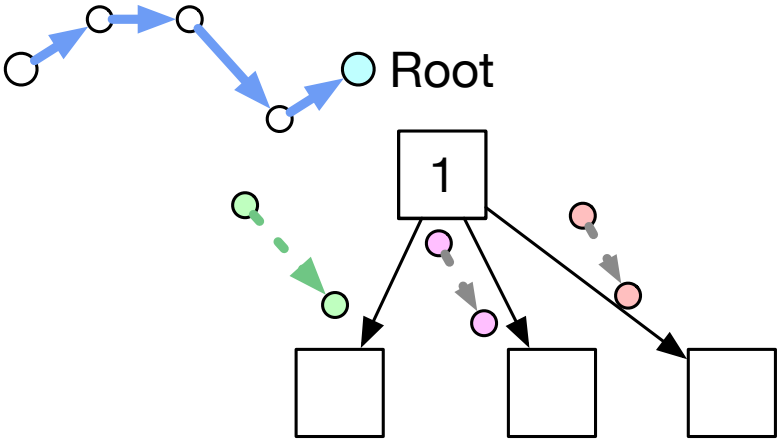
Copy data structure \Leftrightarrow doing *NOTHING* (only keep the old version pointer), $O(1)$

Our Algorithm (1/2): Deviation Preprocessing



BFS on suffix tree and build up heaps $\Rightarrow O(n \log n)$

Our Algorithm (2/2): Efficient Path Searching

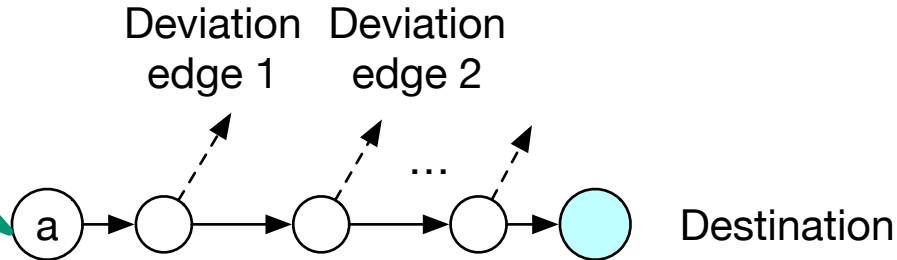


Overall: $O(n \log n + k \log k)$
 = $O(k \log k)$ \nearrow \uparrow
 Preprocessing: $O(n \log n)$

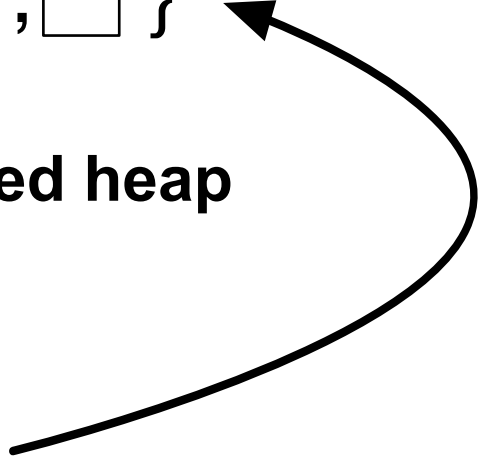
Heap of unexplored nodes $\{ \square, \square, \square \}$

Each exploration: $O(n)$ deviation edges, **but only 1 preprocessed heap**

Heap(a) = {edge 1, edge 2, ...}



Merge!
 $O(\log k)$



Experimental Results

- Implementation on OpenTimer
- 2.1GHz Intel Xeon & 512GB mem
- Compared with:
 - OpenTimer
 - CPU version of the *suffix forest* algorithm [Guo, ICCAD'21]
- On: TAU 2015 benchmarks

On benchmark `leon2`: 4M pins Time (ms)

$k=$	100	100,000	1,000,000
Ours	1481	5476	19472
OpenTimer	4469	218463	1009337
Speed-up	3x	39x	51x
Suffix forest	1441	6834	34572
Speed-up	0.97x	1.24x	1.77x

Conclusions and Future Work

- Near linear-time $O(n \log n + k \log k)$ novel k -shortest path searching algorithm
- **persistent mergeable heaps** and **deviation preprocessing** step to precompute path deviations and speed up path searching.
- **1.7~50x** faster than OpenTimer and other baselines.
- Path constraints
- GPU acceleration
- Common path pessimism removal (CPPR)



Thanks!
Questions are welcome

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